

6.6 控制流: 布尔表达式

- ❖ 布尔表达式通过布尔运算符将布尔变量或关系表达式结合在一起。
- ❖ 布尔表达式的文法:

$$B \rightarrow B \parallel B \mid B \ \&\& \ B \mid !B \mid (B) \mid E \ \text{rel} \ E \mid \text{true} \mid \text{false}$$
$$\text{rel.op} \in \{<, \leq, =, \neq, >, \geq\}$$

布尔表达式

❖ 假设:

- `||`, `&&` 左结合
- 从低到高的优先顺序为: `||`, `&&`, `!`

❖ 翻译时采用短路代码（短路逻辑）:

- `A or B`: 若A为真, 则B无需计算
- `A and B`: 若A为假, 则B无需计算

`! a || b && c || d`

什么情况下, 无需计算c?
什么情况下, 无需计算d?

短路代码

❖ 布尔运算符翻译成跳转指令:

➤ 例: 翻译 `if(x<100 || x >200 && x!=y) x=0;`

```
if x<100 goto L2
```

```
ifFalse x>200 goto L1
```

```
ifFalse x!=y goto L1
```

```
L2: x=0
```

```
L1:
```

用短路代码翻译:

```
if(x<100 && m >200 || x!=y) x=0;
```

控制流翻译

❖ 文法:

$$\begin{aligned} S \rightarrow & \text{if } B \text{ then } S_1 \\ & | \text{if } B \text{ then } S_1 \text{ else } S_2 \\ & | \text{while } B \text{ do } S_1 \end{aligned}$$

❖ 布尔表达式 B 附加两个属性:

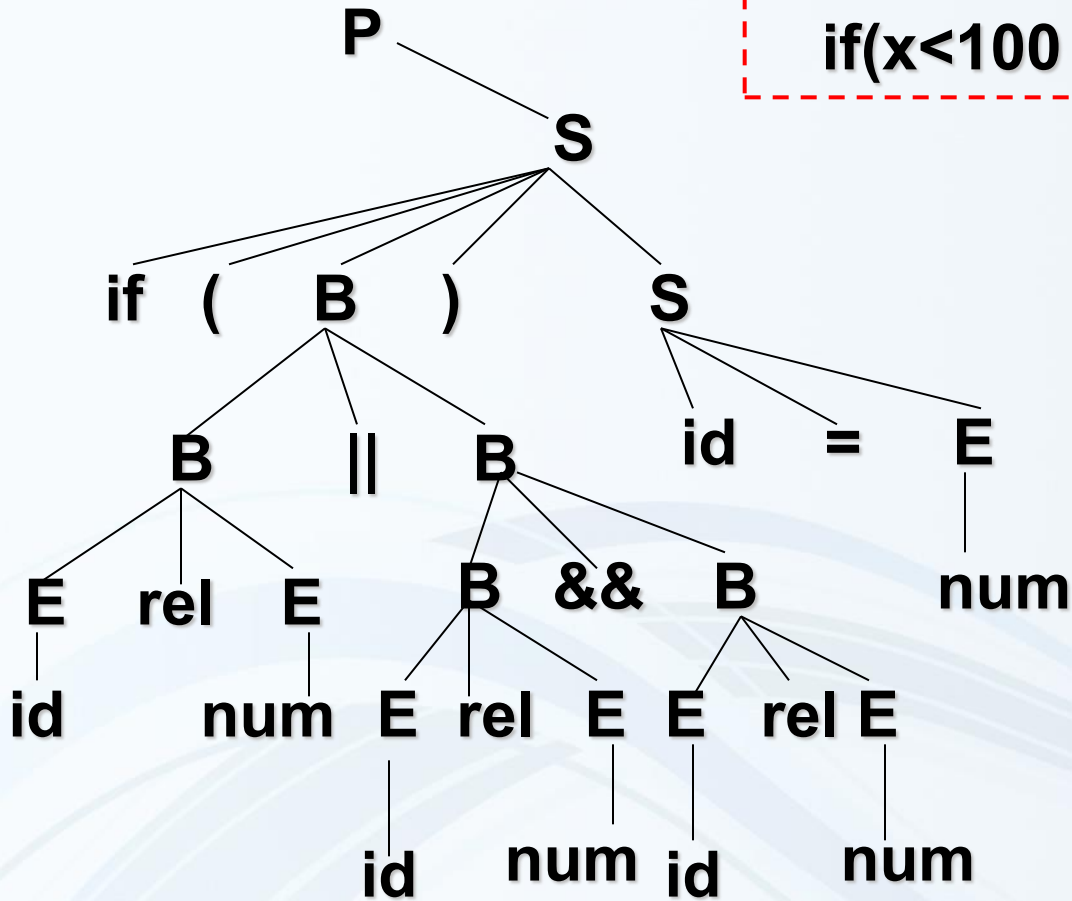
$B.true$: B 为真时的控制流

$B.false$: B 为假时的控制流

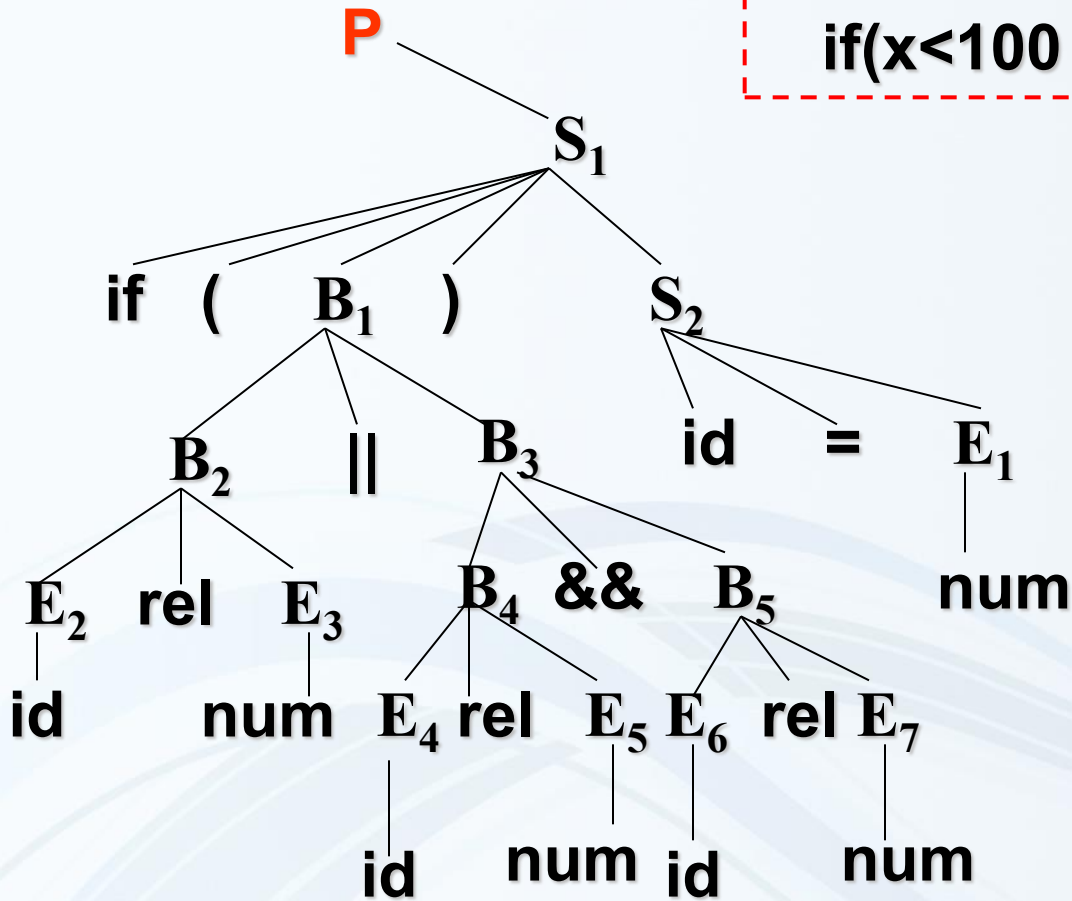
❖ $S.next$ 指向紧跟在 S 后执行的第一条三地址指令

产生式	语义规则
$B \rightarrow B_1 \parallel B_2$	$B_1.true=B_2.true=B.true; B_1.false=newlabel();$ $B_2.false=B.false; B.code=B_1.code \parallel label(B_1.false) \parallel B_2.code$
$B \rightarrow B_1 \ \&\& \ B_2$	$B_1.false=B_2.false=B.false; B_1.true=newlabel();$ $B_2.true=B.true; B.code=B_1.code \parallel label(B_1.true) \parallel B_2.code$
$B \rightarrow !B_1$	$B_1.true=B.false; B_1.false=B.true; B.code=B_1.code;$
$B \rightarrow E_1 \ \text{rel} \ E_2$	$B.code=E_1.code \parallel E_2.code$ $\parallel \text{gen}(\text{'if' } E_1.addr \ \text{rel.op} \ E_2.addr \ \text{'goto' } B.true)$ $\parallel \text{gen}(\text{'goto' } B.false)$
$B \rightarrow \text{true}$	$B.code = \text{gen}(\text{'goto' } B.true)$
$B \rightarrow \text{false}$	$B.code = \text{gen}(\text{'goto' } B.false)$

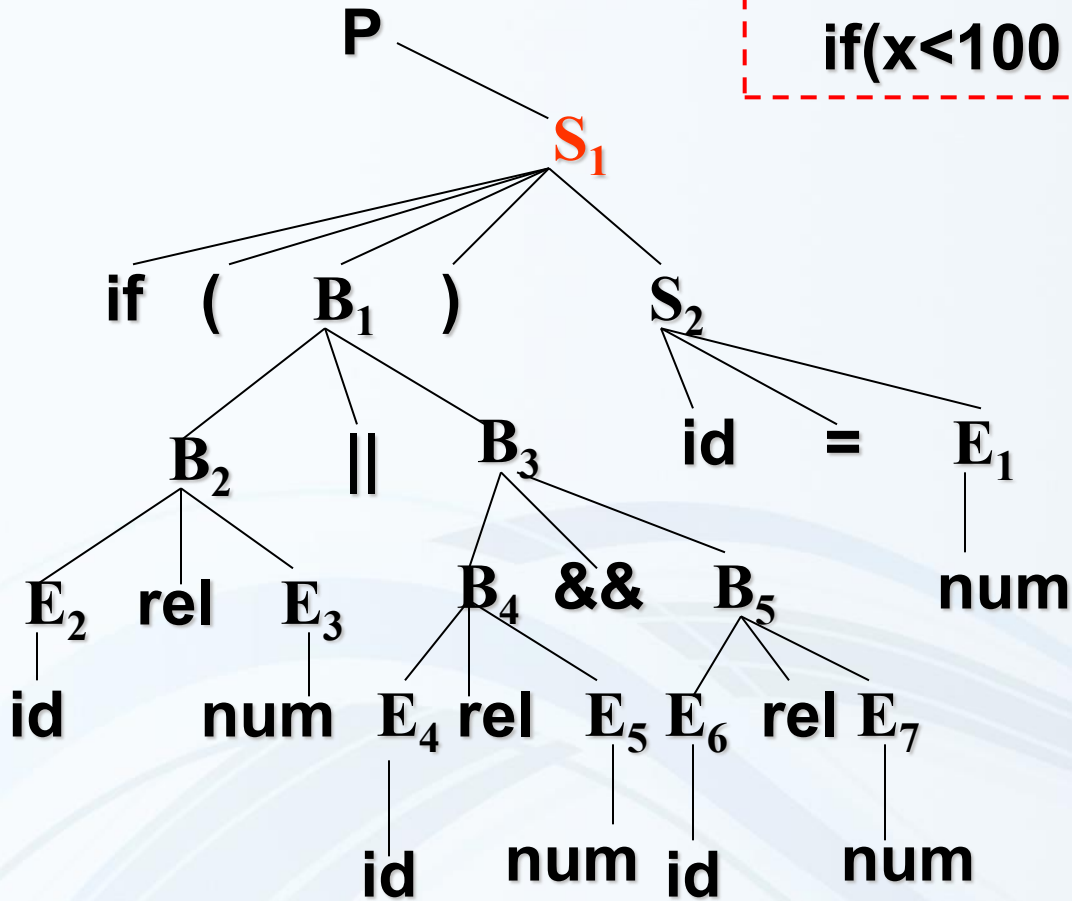
`if(x<100 || x>200 && x!=y) x=0;`



`if(x<100 || x>200 && x!=y) x=0;`

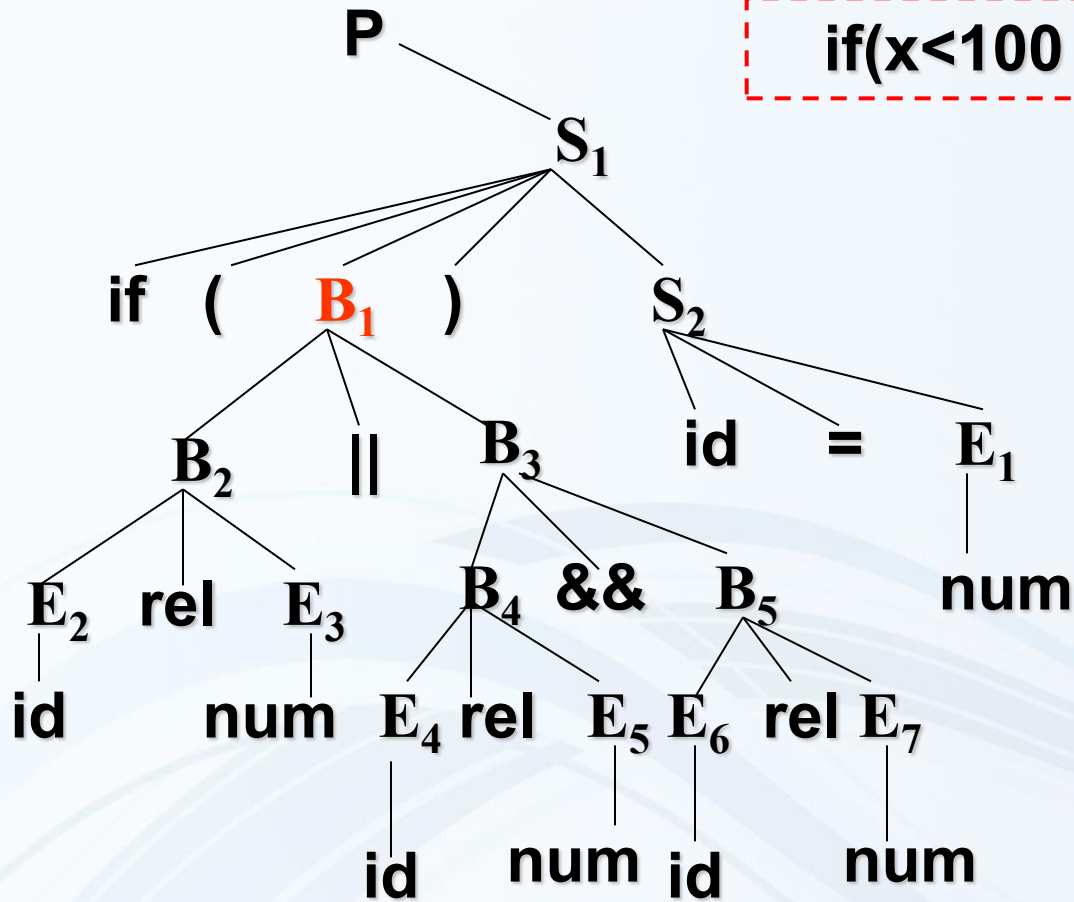


`if(x<100 || x>200 && x!=y) x=0;`



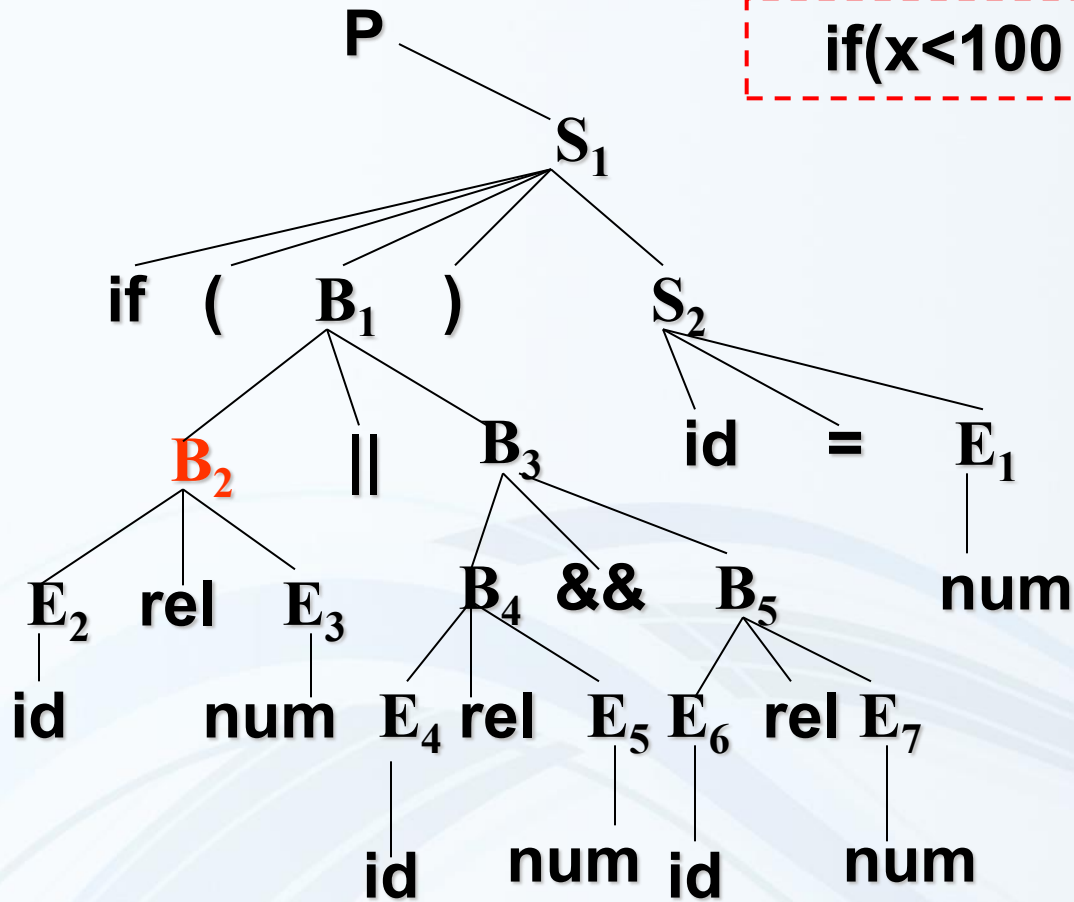
$S_1.next = L_1;$

`if(x<100 || x>200 && x!=y) x=0;`



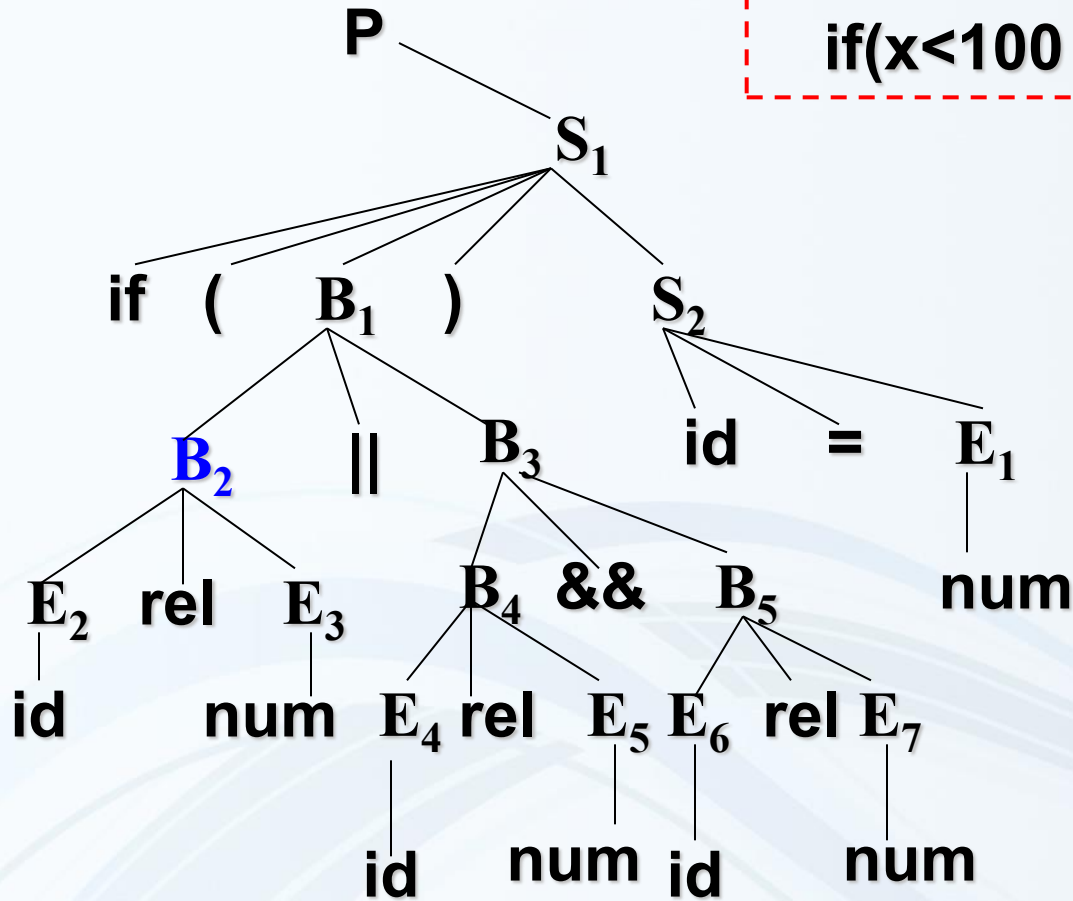
$S_1.next=L_1$; $B_1.true=L_2$; $B_1.false=L_1$;

`if(x<100 || x>200 && x!=y) x=0;`



$S_1.next=L_1; B_1.true=L_2; B_1.false=L_1; B_2.true=L_2; B_2.false=L_3;$

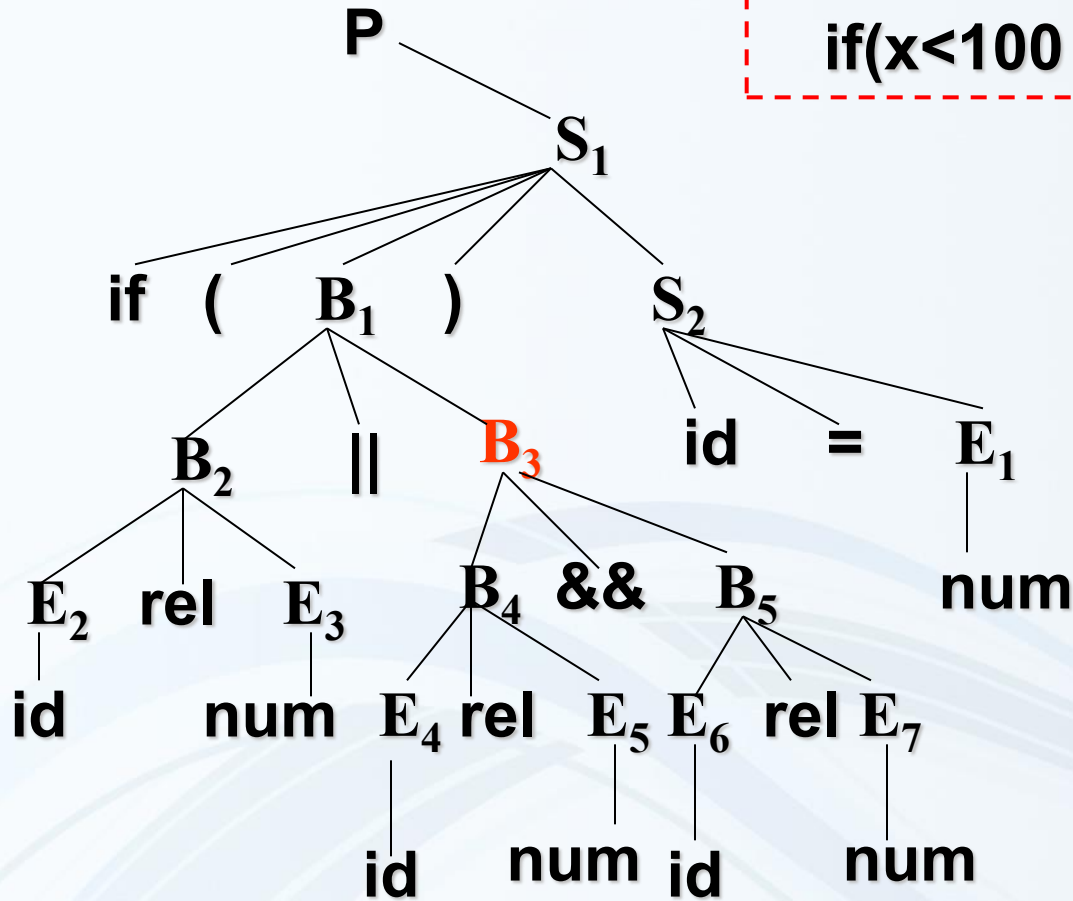
if(x<100 || x>200 && x!=y) x=0;



**if(x<100) goto L2
goto L3**

$S_1.next=L_1; B_1.true=L_2; B_1.false=L_1; B_2.true=L_2; B_2.false=L_3;$

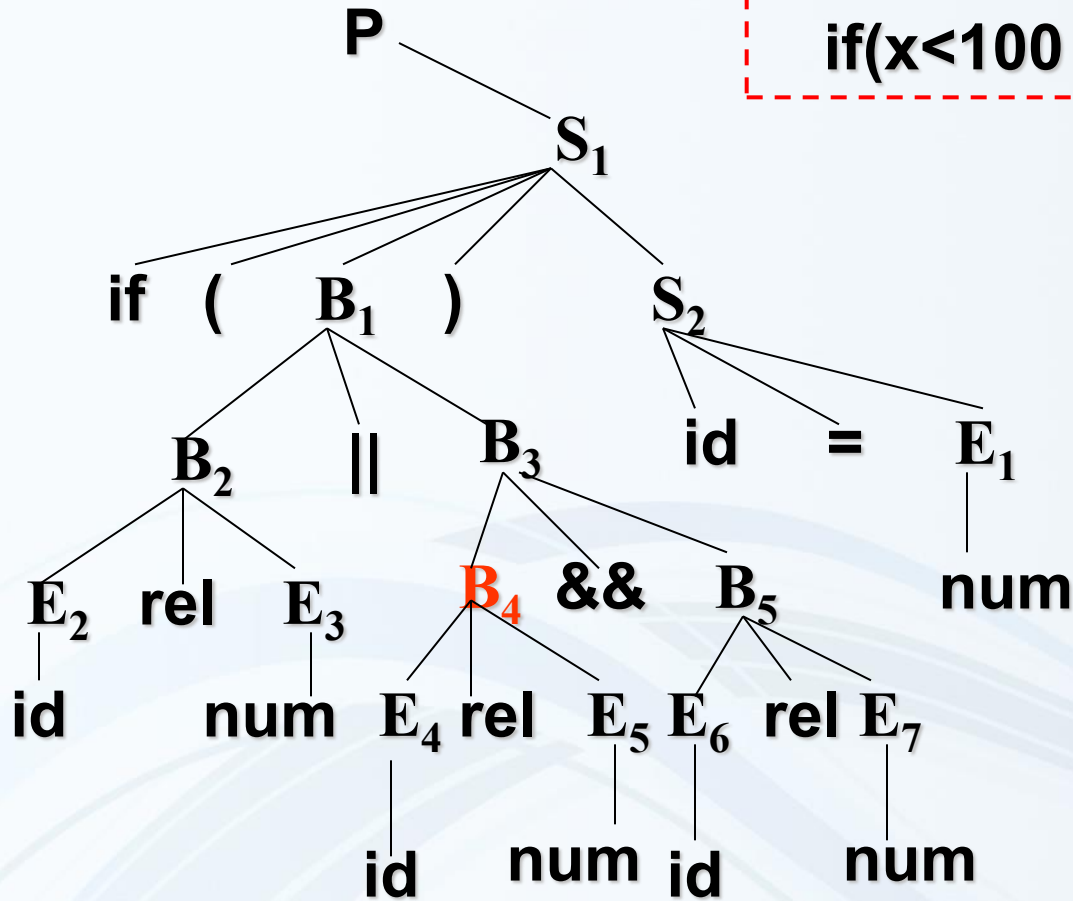
if(x<100 || x>200 && x!=y) x=0;



if(x<100) goto L2
goto L3

$S_1.next=L_1; B_1.true=L_2; B_1.false=L_1; B_2.true=L_2; B_2.false=L_3; B_3.true=L_2; B_3.false=L_1;$

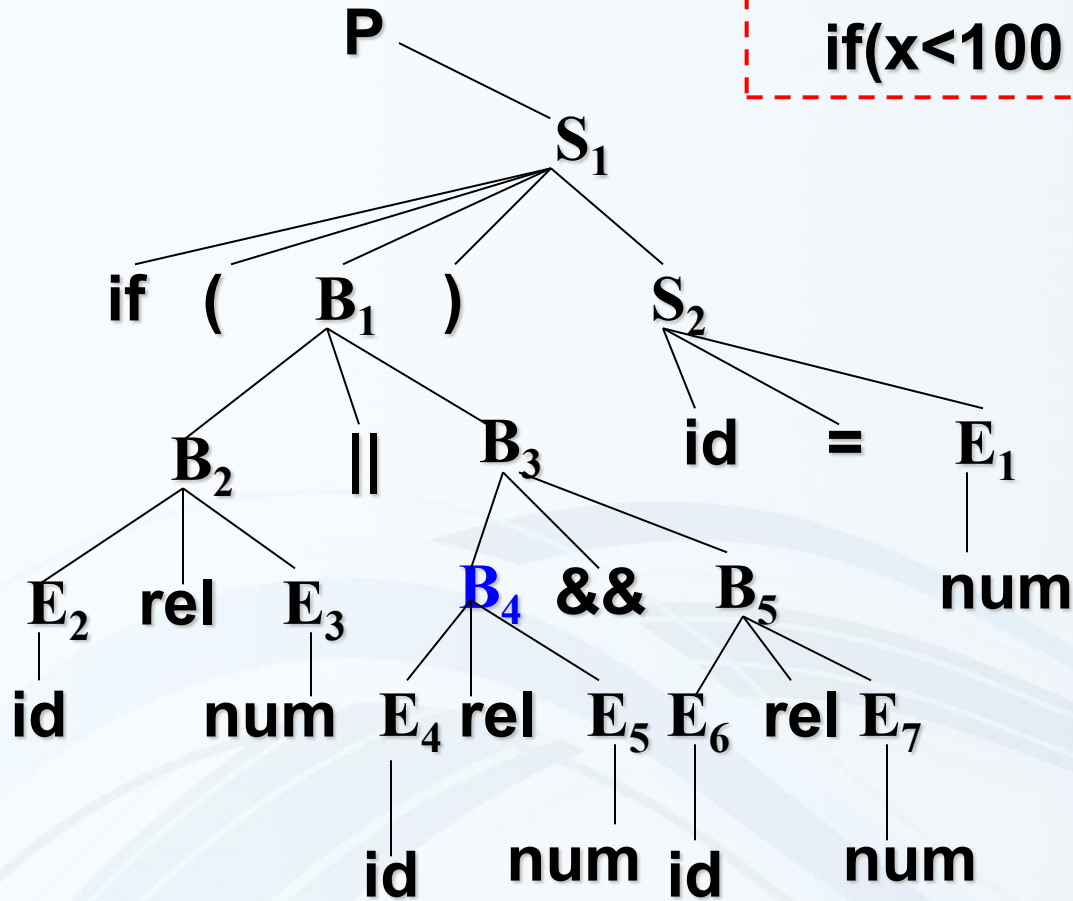
if(x<100 || x>200 && x!=y) x=0;



if(x<100) goto L2
goto L3

$S_1.next=L_1; B_1.true=L_2; B_1.false=L_1; B_2.true=L_2; B_2.false=L_3; B_3.true=L_2; B_3.false=L_1; B_4.true=L_4; B_4.false=L_1;$

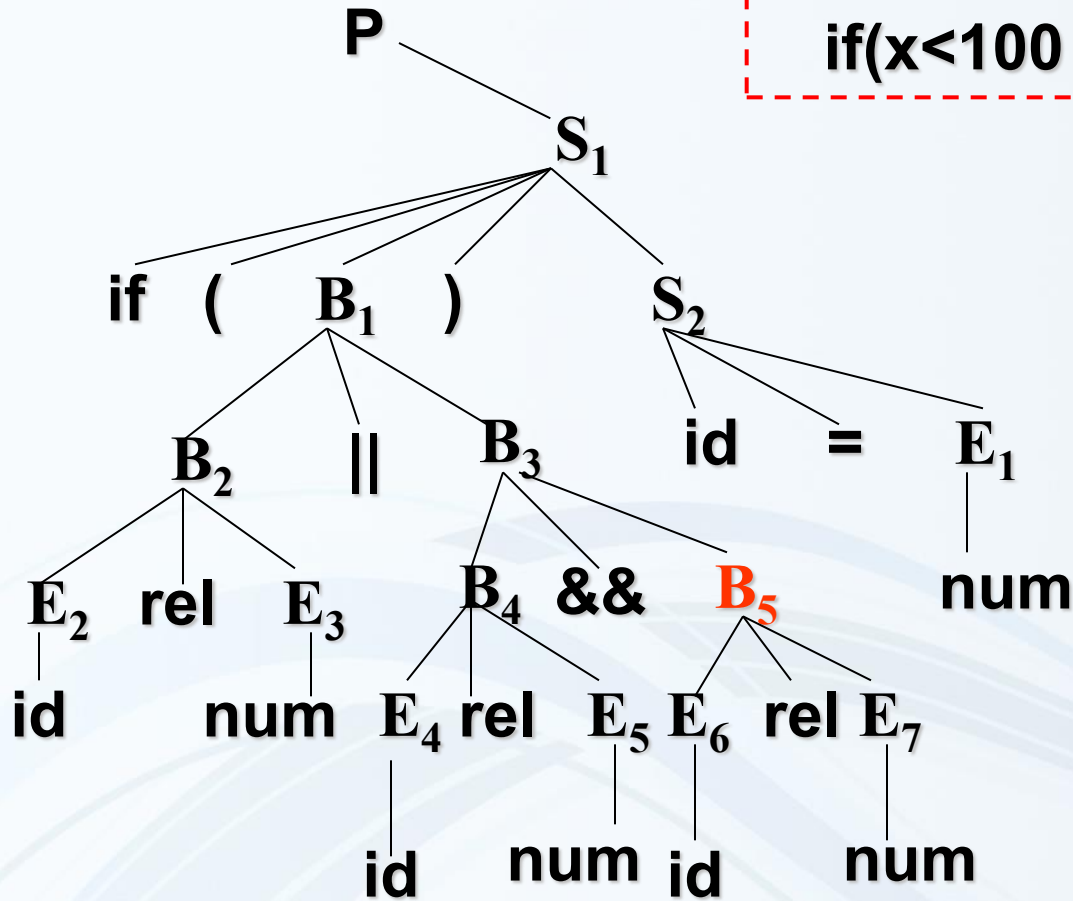
if(x<100 || x>200 && x!=y) x=0;



if(x<100) goto L2
goto L3
if(x>200) goto L4
goto L1

$S_1.next=L_1; B_1.true=L_2; B_1.false=L_1; B_2.true=L_2; B_2.false=L_3; B_3.true=L_2; B_3.false=L_1; B_4.true=L_4; B_4.false=L_1;$

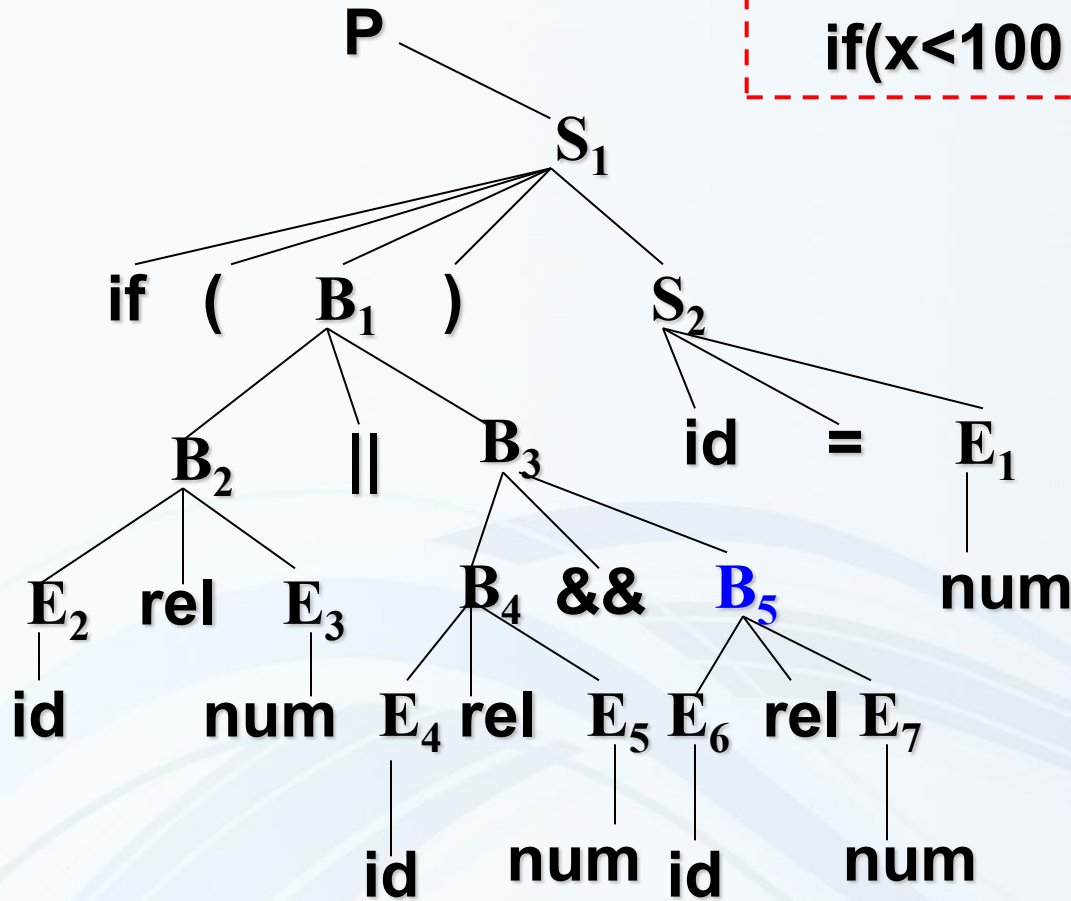
if(x<100 || x>200 && x!=y) x=0;



if(x<100) goto L2
 goto L3
 if(x>200) goto L4
 goto L1

$S_1.next=L_1; B_1.true=L_2; B_1.false=L_1; B_2.true=L_2; B_2.false=L_3; B_3.true=L_2; B_3.false=L_1; B_4.true=L_4; B_4.false=L_1; B_5.true=L_2; B_5.false=L_1;$

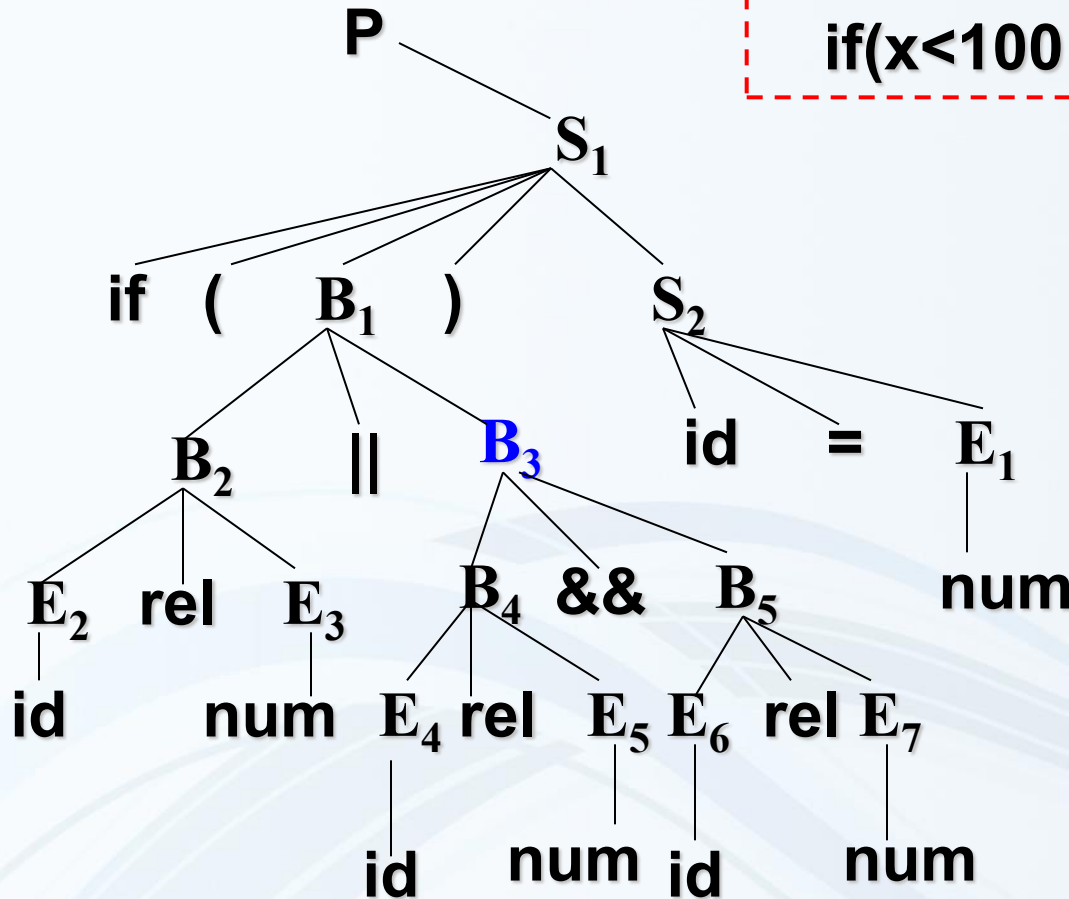
if(x<100 || x>200 && x!=y) x=0;



if(x<100) goto L2
 goto L3
 if(x>200) goto L4
 goto L1
if(x!=y) goto L2
goto L1

$S_1.next=L_1; B_1.true=L_2; B_1.false=L_1; B_2.true=L_2; B_2.false=L_3; B_3.true=L_2; B_3.false=L_1; B_4.true=L_4; B_4.false=L_1; B_5.true=L_2; B_5.false=L_1;$

if(x<100 || x>200 && x!=y) x=0;



if(x<100) goto L2

goto L3

if(x>200) goto L4

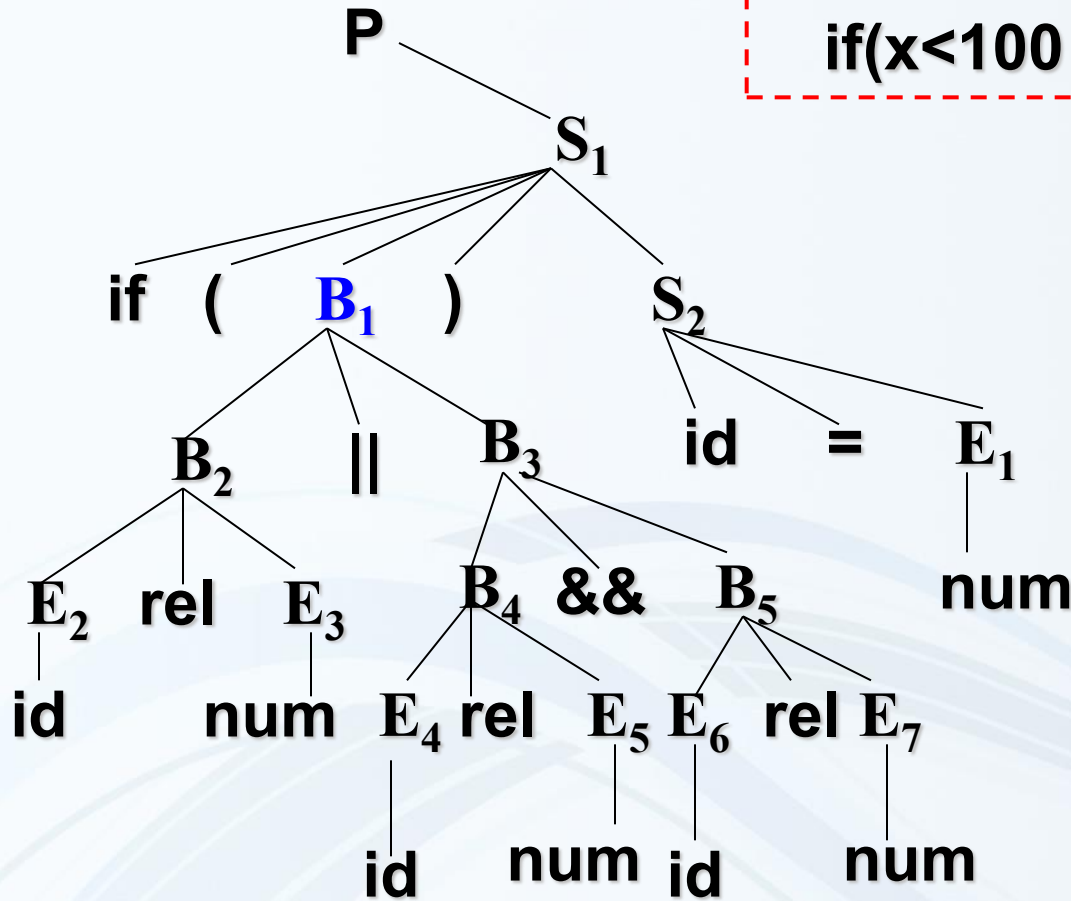
goto L1

L4:if(x!=y) goto L2

goto L1

$S_1.next=L_1; B_1.true=L_2; B_1.false=L_1; B_2.true=L_2; B_2.false=L_3; B_3.true=L_2; B_3.false=L_1; B_4.true=L_4; B_4.false=L_1; B_5.true=L_2; B_5.false=L_1;$

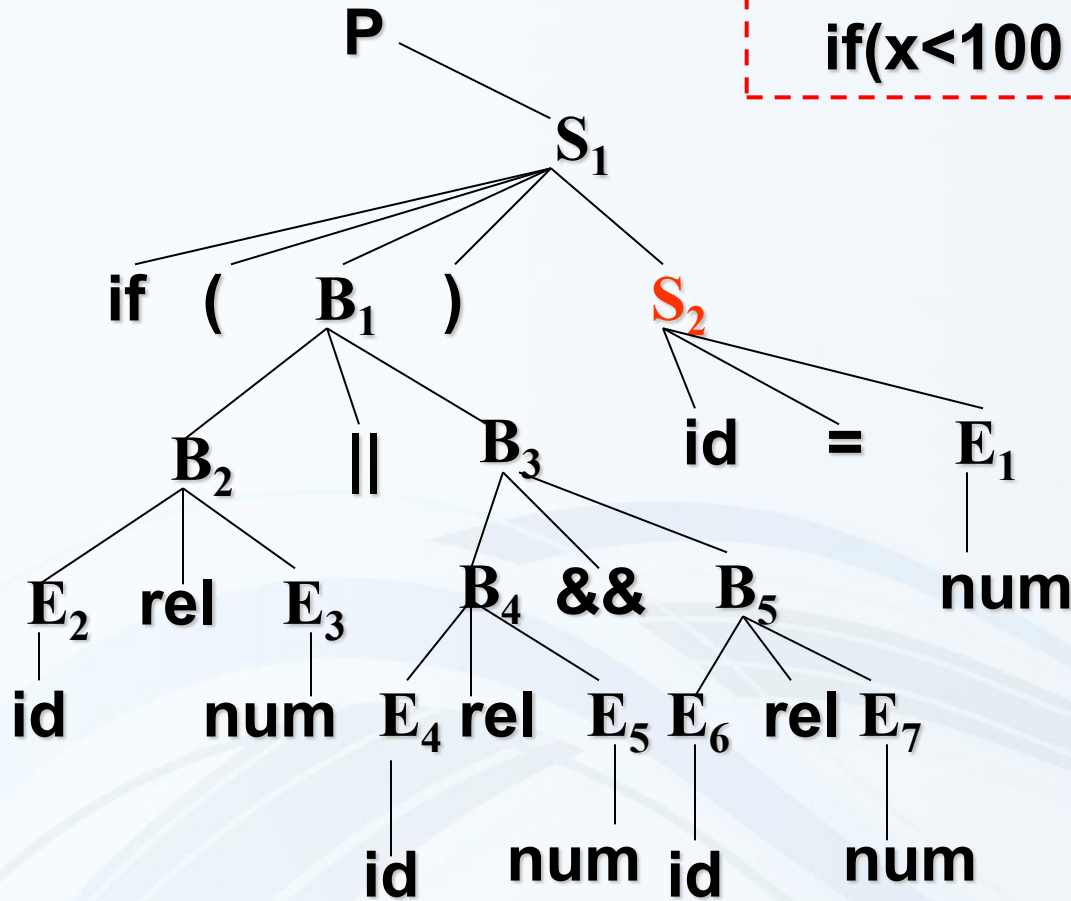
if(x<100 || x>200 && x!=y) x=0;



if(x<100) goto L2
 goto L3
L3:if(x>200) goto L4
 goto L1
 L4:if(x!=y) goto L2
 goto L1

$S_1.next=L_1; B_1.true=L_2; B_1.false=L_1; B_2.true=L_2; B_2.false=L_3; B_3.true=L_2; B_3.false=L_1; B_4.true=L_4; B_4.false=L_1; B_5.true=L_2; B_5.false=L_1;$

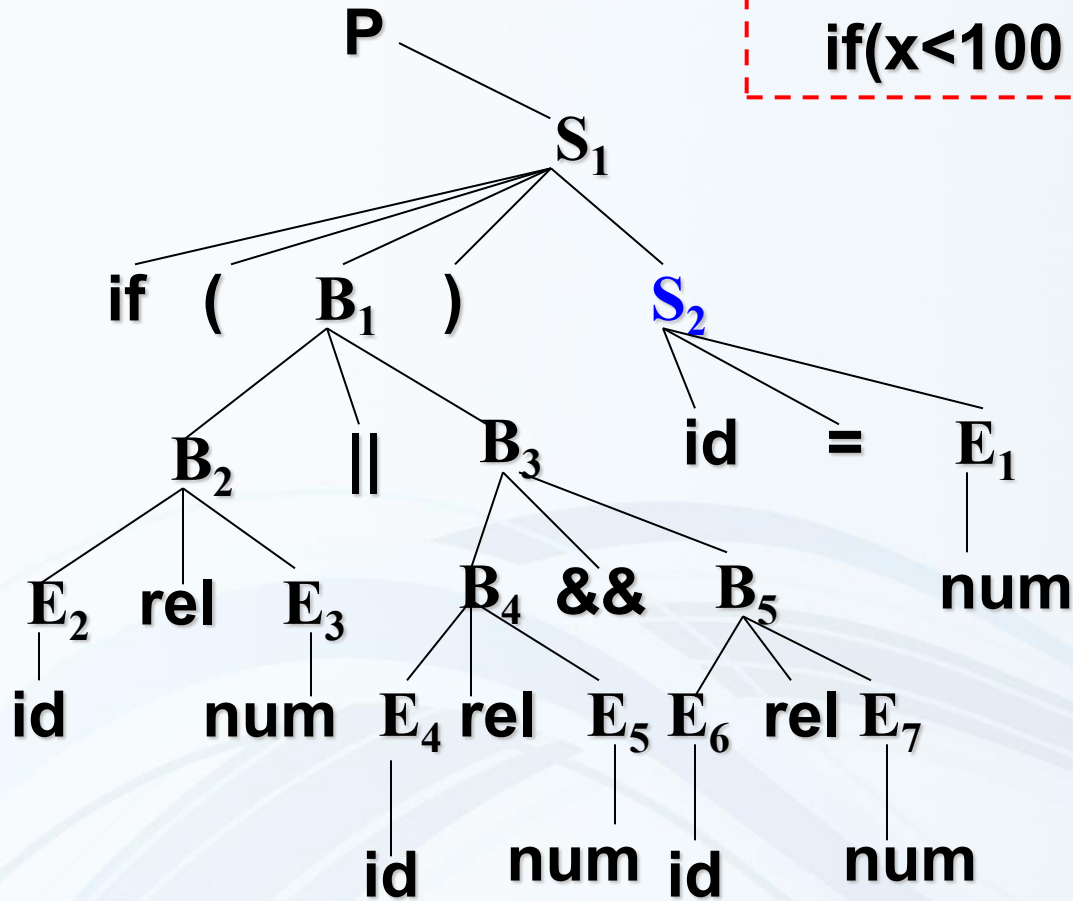
if(x<100 || x>200 && x!=y) x=0;



if(x<100) goto L2
 goto L3
 L3:if(x>200) goto L4
 goto L1
 L4:if(x!=y) goto L2
 goto L1

$S_1.next=L_1; B_1.true=L_2; B_1.false=L_1; B_2.true=L_2; B_2.false=L_3; B_3.true=L_2; B_3.false=L_1; B_4.true=L_4; B_4.false=L_1; B_5.true=L_2; B_5.false=L_1; S_2.next=L_1;$

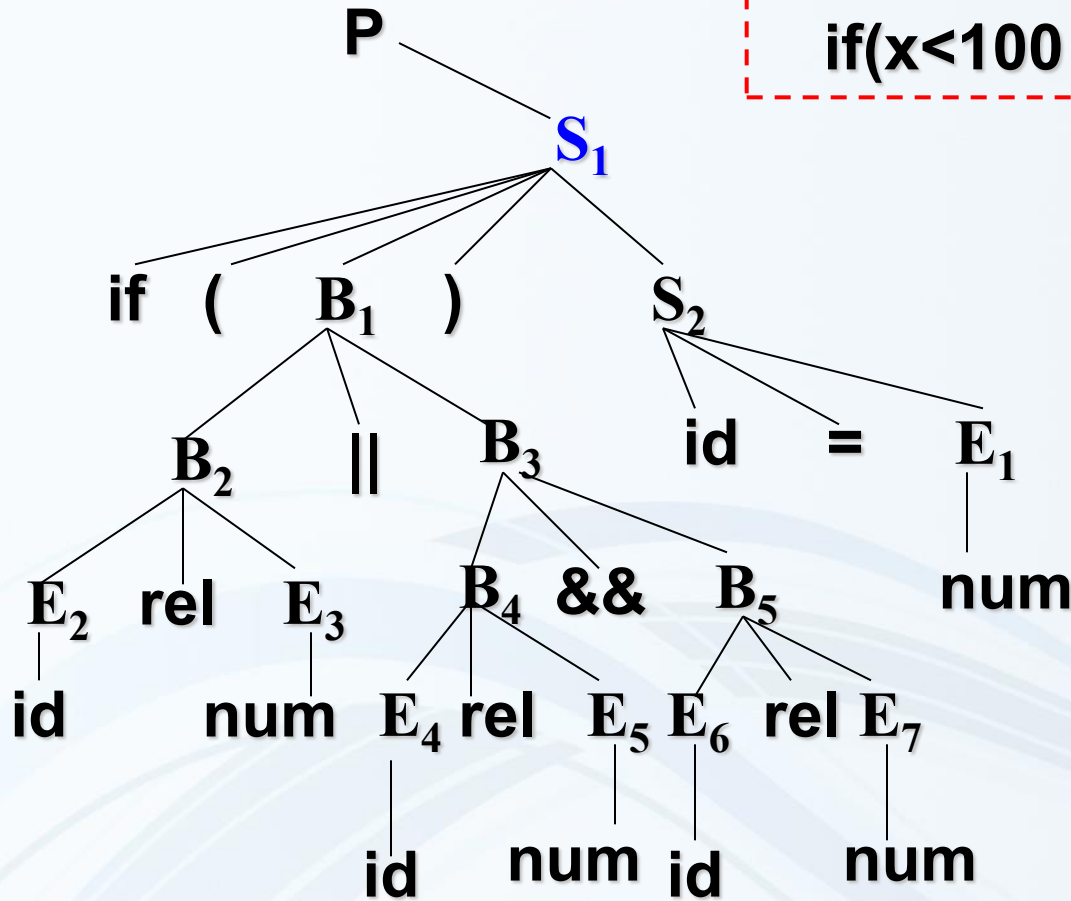
if(x<100 || x>200 && x!=y) x=0;



if(x<100) goto L2
 goto L3
 L3:if(x>200) goto L4
 goto L1
 L4:if(x!=y) goto L2
 goto L1
x=0

$S_1.next=L_1; B_1.true=L_2; B_1.false=L_1; B_2.true=L_2; B_2.false=L_3; B_3.true=L_2; B_3.false=L_1; B_4.true=L_4; B_4.false=L_1; B_5.true=L_2; B_5.false=L_1; S_2.next=L_1;$

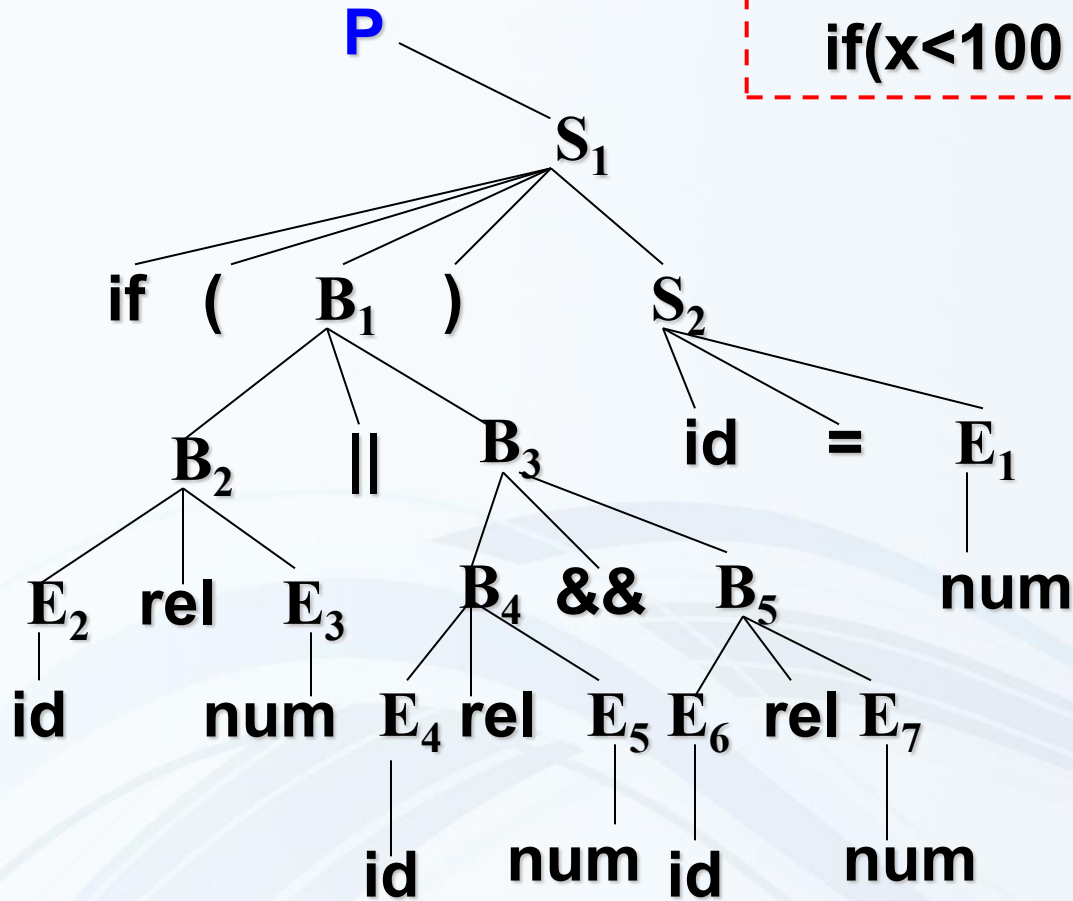
if(x<100 || x>200 && x!=y) x=0;



if(x<100) goto L2
 goto L3
 L3:if(x>200) goto L4
 goto L1
 L4:if(x!=y) goto L2
 goto L1
L2:x=0

$S_1.next=L_1; B_1.true=L_2; B_1.false=L_1; B_2.true=L_2; B_2.false=L_3; B_3.true=L_2; B_3.false=L_1; B_4.true=L_4; B_4.false=L_1; B_5.true=L_2; B_5.false=L_1; S_2.next=L_1;$

if(x<100 || x>200 && x!=y) x=0;



if(x<100) goto L2
 goto L3
 L3:if(x>200) goto L4
 goto L1
 L4:if(x!=y) goto L2
 goto L1
 L2:x=0

L1:

$S_1.next=L_1; B_1.true=L_2; B_1.false=L_1; B_2.true=L_2; B_2.false=L_3; B_3.true=L_2; B_3.false=L_1; B_4.true=L_4; B_4.false=L_1; B_5.true=L_2; B_5.false=L_1; S_2.next=L_1;$

```
if(x<100 || x>200 && x!=y) x=0;
```

```
100: if (x<100) goto
```

```
101: goto
```

```
102: if(x>200) goto
```

```
103: goto
```

```
104: if(x!=y) goto
```

```
105: goto
```

```
106: x=0
```

```
107:
```

```
if(x<100) goto L2
```

```
goto L3
```

```
L3:if(x>200) goto L4
```

```
goto L1
```

```
L4:if(x!=y) goto L2
```

```
goto L1
```

```
L2:x=0
```

```
L1:
```

控制流翻译的SDD

产生式	语义规则
$P \rightarrow S$	$S.next = newlabel(); P.code = S.code \parallel label(S.next)$
$S \rightarrow assign$	$S.code = assign.code$
$S \rightarrow if(B) S_1$	$B.true = newlabel(); B.false = S_1.next = S.next;$ $S.code = B.code \parallel label(B.true) \parallel S_1.code$
$S \rightarrow if(B) S_1$ $else S_2$	$B.true = newlabel(); B.false = newlabel();$ $S_1.next = S_2.next = S.next;$ $S.code = B.code \parallel label(B.true) \parallel S_1.code$ $\parallel gen('goto' S.next) \parallel label(B.false) \parallel S_2.code$

控制流翻译的SDD(II)

产生式	语义规则
$S \rightarrow \text{while}(B)S_1$	$\text{begin}=\text{newlabel}()$; $B.\text{true}=\text{newlabel}()$; $B.\text{false}=S.\text{next}$; $S_1.\text{next}=\text{begin}$; $S.\text{code}=\text{label}(\text{begin}) \parallel B.\text{code} \parallel \text{label}(B.\text{true}) \parallel$ $S_1.\text{code} \parallel \text{gen}(\text{'goto' begin})$
$S \rightarrow S_1 S_2$	$S_1.\text{next}=\text{newlabel}()$; $S_2.\text{next}=S.\text{next}$; $S.\text{code} = S_1.\text{code} \parallel \text{label}(S_1.\text{next}) \parallel S_2.\text{code}$

产生式 $S \rightarrow \text{repeat } S_1 \text{ until } B$ 对应的语义规则是什么?

6.7 回填 *Backpatching*

❖ 两趟式执行语法制导定义（L属性定义）：

1. 将标记作为继承属性传至生成跳转指令处
2. 绑定标记和地址

❖ 一趟式执行：（S属性定义）：

- 用综合属性 *truelist* 和 *falselist* 管理跳转指令中的标记
- 采用回填：在目标地址确定时，填充这些标记

回填(II)

❖ 三个函数:

1. *makelist*(**i**) 建立新表,表中加入元素**i**(**i**是指令数组的下标),返回表指针
2. *merge*(p_1, p_2) 合并 p_1 、 p_2 指向的表,返回 p_2
3. *backpatch*(p, i) 将**i回填**入 p 指向列表中每一语句

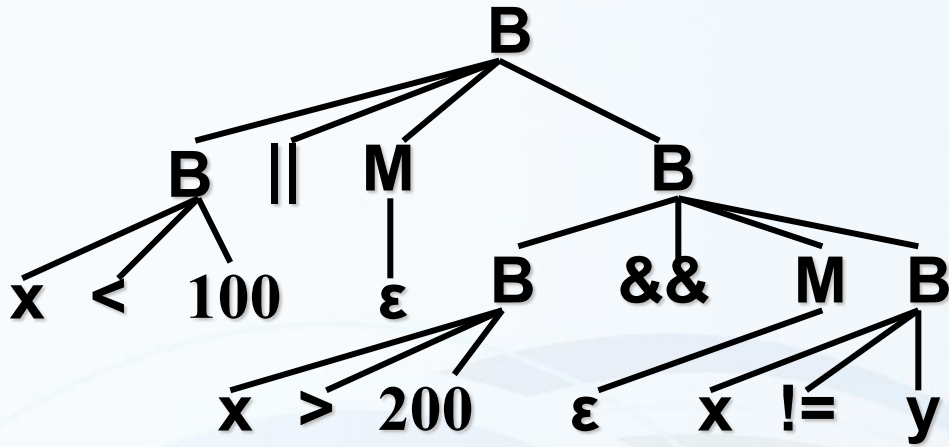
❖ **nextinstr**: 下一条指令的地址

布尔表达式的回填(II)

$B \rightarrow E_1 \text{ rel } E_2$	<pre>{B.truelist = makelist(nextinstr); B.falselist = makelist(nextinstr+1); gen('if' E₁.addr rel.op E₂.addr 'goto _'); gen('goto _');}</pre>
$B \rightarrow \text{true}$	<pre>{B.truelist = makelist(nextinstr); gen('goto _');}</pre>
$B \rightarrow \text{false}$	<pre>{B.falselist = makelist(nextinstr); gen('goto _');}</pre>

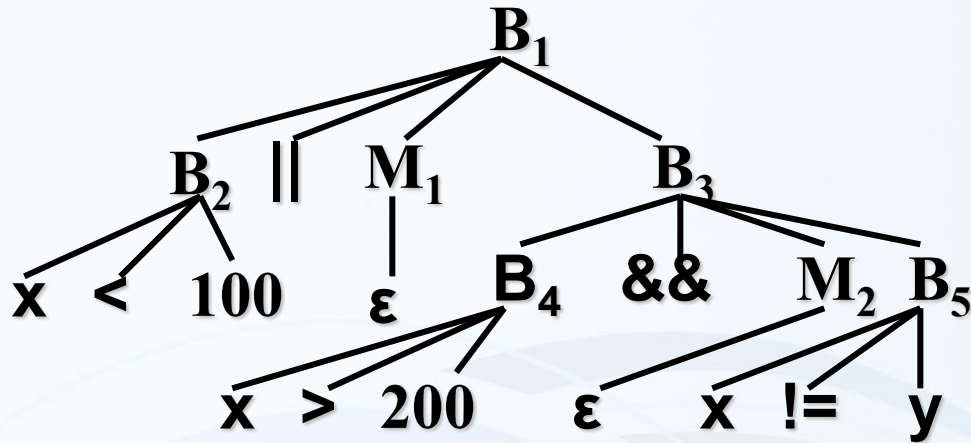
为什么 $B.falselist = makelist(\text{nextinstr}+1);$?

布尔表达式的回填(例)

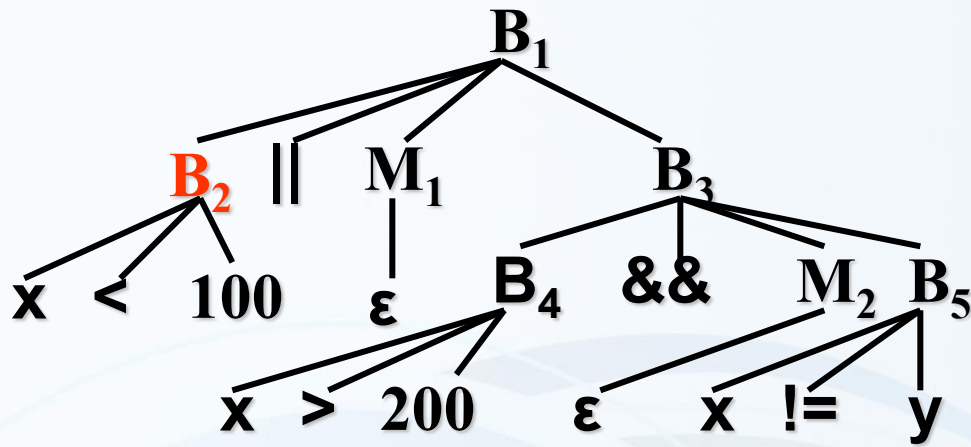


$x < 100 \ || \ x > 200 \ \&\& \ x \neq y$

布尔表达式的回填(例)



布尔表达式的回填(例)



100:if x < 100 goto

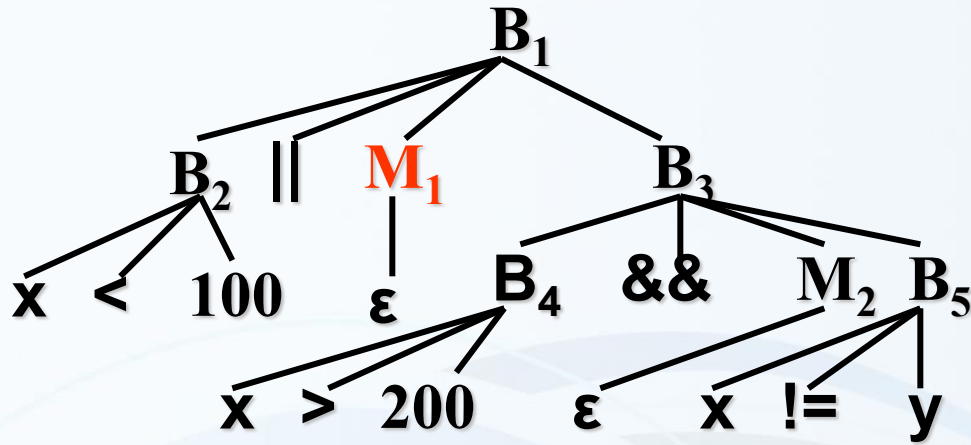
101:goto

$B_2.t = \{100\}; B_2.f = \{101\};$

布尔表达式的回填(例)

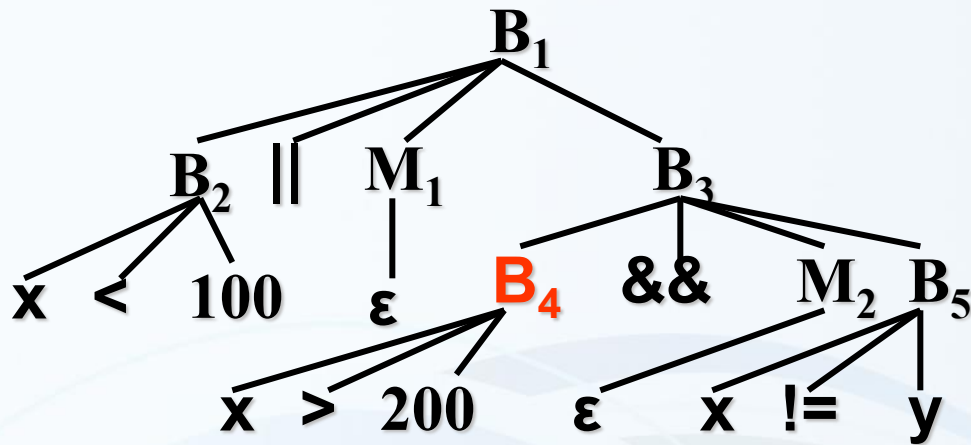
100:if x < 100 goto

101:goto



$B_2.t = \{100\}$; $B_2.f = \{101\}$; $M_1.i = 102$;

布尔表达式的回填(例)



100:if x < 100 goto

101:goto

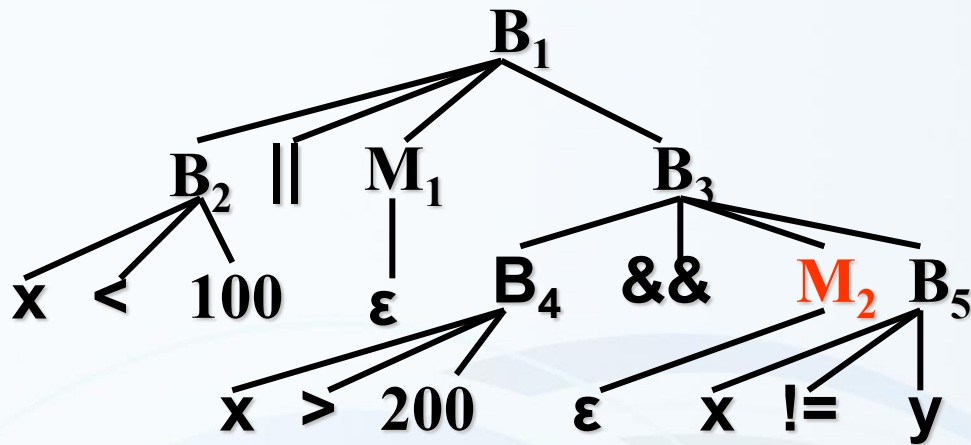
102:if x > 200 goto

103:goto

$B_2.t = \{100\}; B_2.f = \{101\}; M_1.i = 102;$

$B_4.t = \{102\}; B_4.f = \{103\};$

布尔表达式的回填(例)



100:if x < 100 goto

101:goto

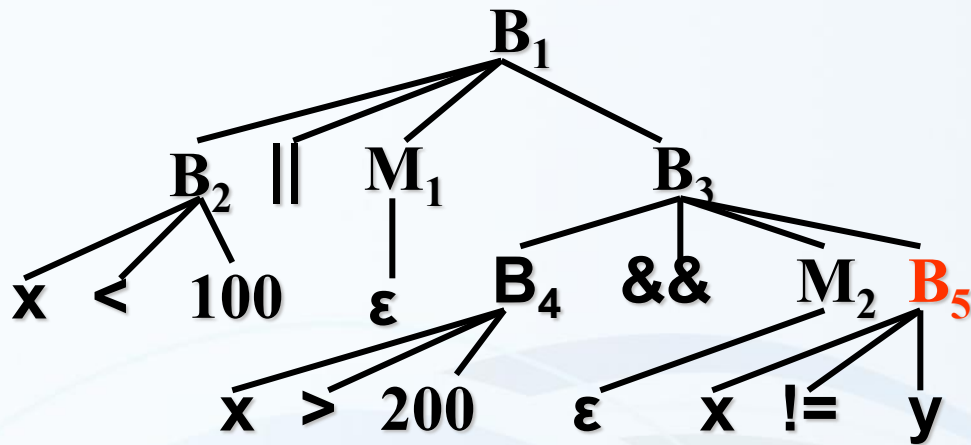
102:if x > 200 goto

103:goto

$B_2.t = \{100\}; B_2.f = \{101\}; M_1.i = 102;$

$B_4.t = \{102\}; B_4.f = \{103\}; M_2.i = 104;$

布尔表达式的回填(例)



100:if x < 100 goto

101:goto

102:if x > 200 goto

103:goto

104:if x!=y goto

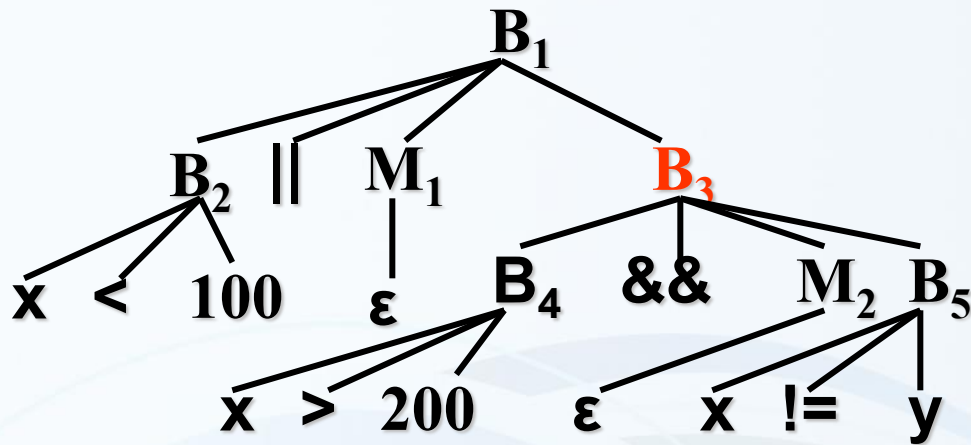
105:goto

$B_2.t = \{100\}; B_2.f = \{101\}; M_1.i = 102;$

$B_4.t = \{102\}; B_4.f = \{103\}; M_2.i = 104;$

$B_5.t = \{104\}; B_5.f = \{105\};$

布尔表达式的回填(例)



100:if x < 100 goto

101:goto

102:if x > 200 goto **104**

103:goto

104:if x!=y goto

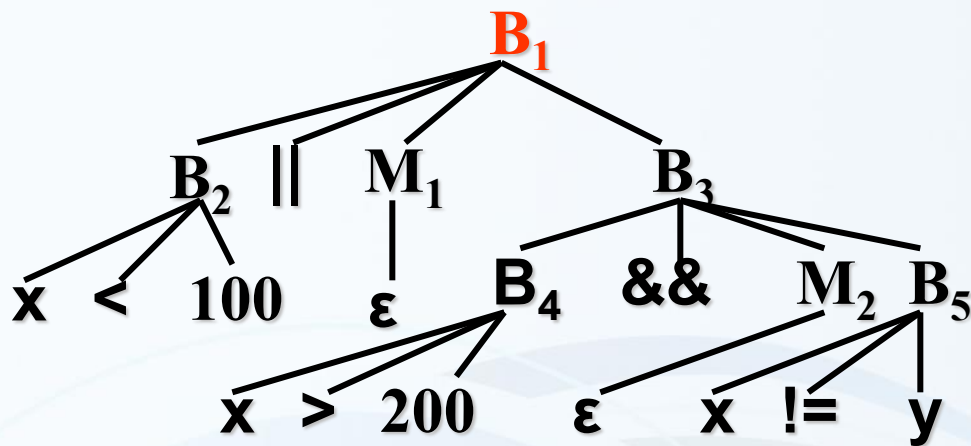
105:goto

$B_2.t = \{100\}; B_2.f = \{101\}; M_1.i = 102;$

$B_4.t = \{102\}; B_4.f = \{103\}; M_2.i = 104;$

$B_5.t = \{104\}; B_5.f = \{105\}; B_3.t = \{104\};$

$B_3.f = \{103, 105\};$



```

100:if x < 100 goto
101:goto 102
102:if x > 200 goto 104
103:goto
104:if x!=y goto
105:goto
  
```

```

B2.t={100}; B2.f={101}; M1.i=102;
B4.t={102}; B4.f={103}; M2.i=104;
B5.t={104}; B5.f={105}; B3.t={104};
B3.f={103, 105}; B1.t={100, 104};
B1.f={103,105}
  
```

真出口填100,104行，假出口填103,105行

控制流语句

- ❖ 语句S带有一个未填充的跳转链 *nextlist*, 该链同样使用回填策略来填充
- ❖ *S.nextlist* 为跳转到S之后的跳转指令列表

```
1) S → if ( B ) M S1 { backpatch(B.truelist, M.instr);  
                        S.nextlist = merge(B.falselist, S1.nextlist); }
```

```
2) S → if ( B ) M1 S1 N else M2 S2  
      { backpatch(B.truelist, M1.instr);  
        backpatch(B.falselist, M2.instr);  
        temp = merge(S1.nextlist, N.nextlist);  
        S.nextlist = merge(temp, S2.nextlist); }
```

```
3) S → while M1 ( B ) M2 S1  
      { backpatch(S1.nextlist, M1.instr);  
        backpatch(B.truelist, M2.instr);  
        S.nextlist = B.falselist;  
        gen('goto' M1.instr); }
```

控制流语句 (II)

- | | |
|-----------------------------|---|
| 4) $S \rightarrow \{ L \}$ | $\{ S.nextlist = L.nextlist; \}$ |
| 5) $S \rightarrow A ;$ | $\{ S.nextlist = \text{null}; \}$ |
| 6) $M \rightarrow \epsilon$ | $\{ M.instr = \text{nextinstr}; \}$ |
| 7) $N \rightarrow \epsilon$ | $\{ N.nextlist = \text{makelist}(\text{nextinstr});$
$\text{gen}(\text{'goto -'}); \}$ |
| 8) $L \rightarrow L_1 M S$ | $\{ \text{backpatch}(L_1.nextlist, M.instr);$
$L.nextlist = S.nextlist; \}$ |
| 9) $L \rightarrow S$ | $\{ L.nextlist = S.nextlist; \}$ |

例 `while (E_1) {`
 `if (E_2)`
 `while (E_3)`
 S_1 ;
 `else {`
 `if (E_4)`
 S_2 ;
 S_3
 `}`
 `}`

i_1 : Code for E_1
 i_2 : Code for E_2
 i_3 : Code for E_3
 i_4 : Code for S_1
 i_5 : Code for E_4
 i_6 : Code for S_2
 i_7 : Code for S_3
 i_8 : ...

求:

- (a) E_3 .false (b) S_2 .next (c) E_4 .false
 (d) S_1 .next (e) E_2 .true

例

```
while ( $E_1$ ) {  
    if ( $\bar{E}_2$ )  
        while ( $E_3$ )  
             $S_1$ ;  
    else {  
        if ( $E_4$ )  
             $S_2$ ;  
         $S_3$   
    }  
}
```

i_1 : Code for E_1
 i_2 : Code for E_2
 i_3 : Code for E_3
 i_4 : Code for S_1
 i_5 : Code for E_4
 i_6 : Code for S_2
 i_7 : Code for S_3
 i_8 : ...

S_4 : while (E_3) S_1

S_5 : if (E_4) S_2 }

S_6 : 包括 S_5 和 S_3 的代码段

S_7 : if(E_2) S_4 else S_6

S_8 : 全部程序

确定: (a) S_4 .next

(b) S_5 .next

(c) S_6 .next

(d) S_7 .next

(e) S_8 .next